

# PAC-Learning & Monitoring

*How machine learning could  
help runtime verification.*

# A high-level overview.

*Focus on problems, not results.*

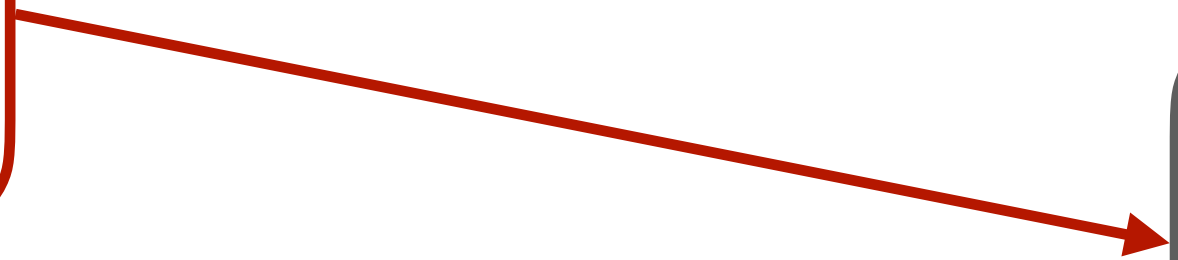
# Formal Verification.

*Proving the correctness of a system.*

# Monitoring.

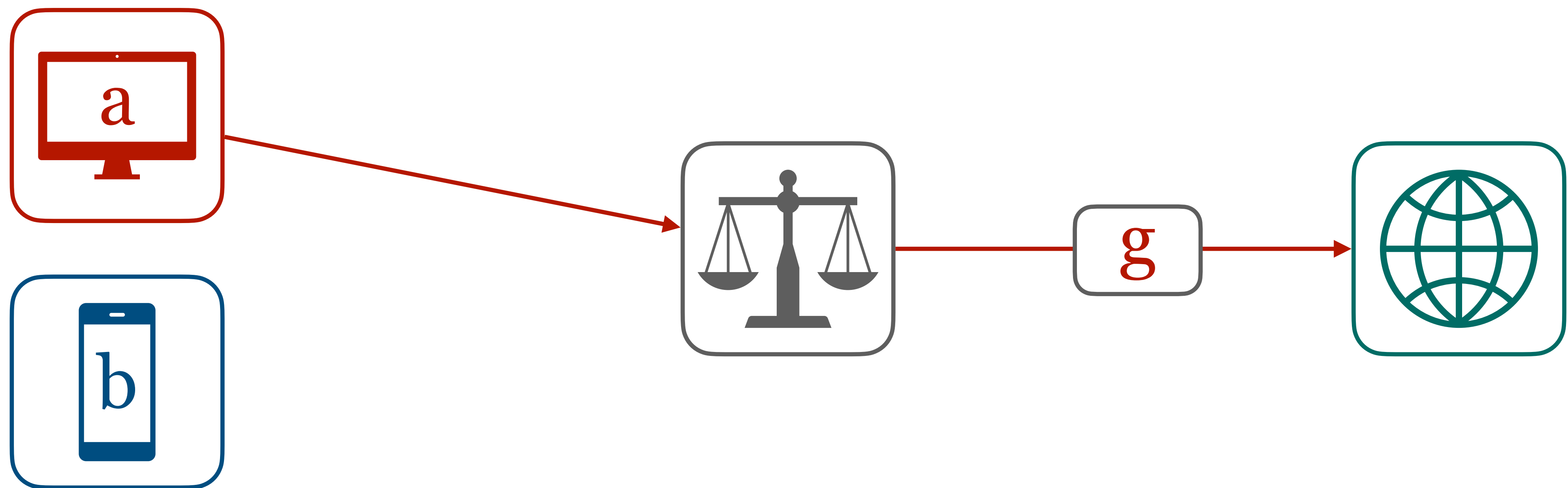
*Proving the correctness of a system  
on a particular run at runtime.*



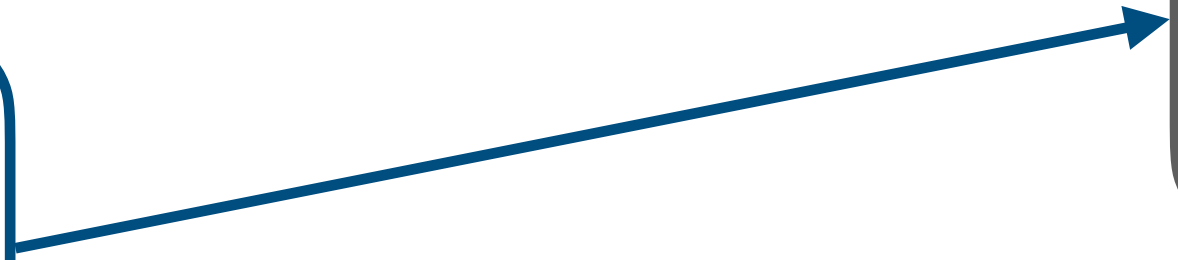
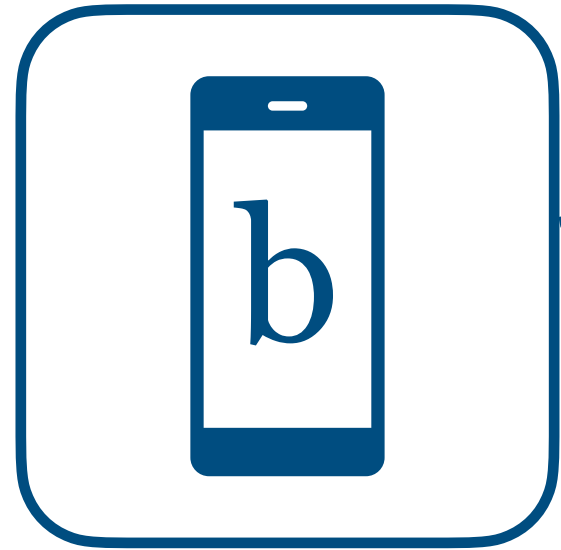


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a



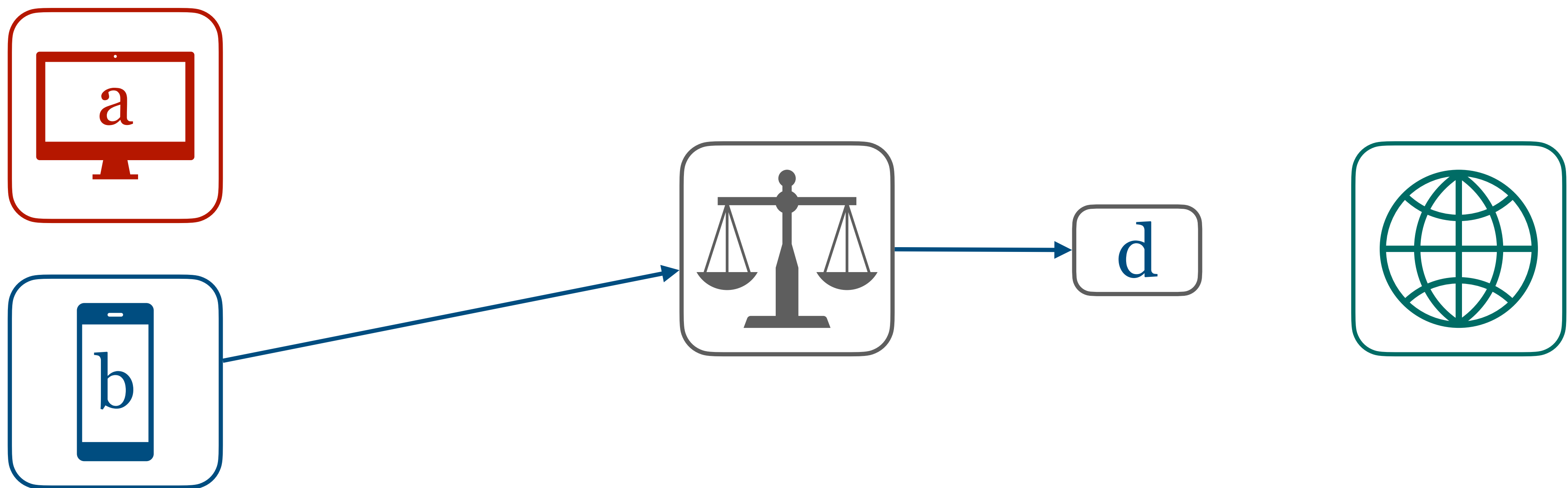
a g



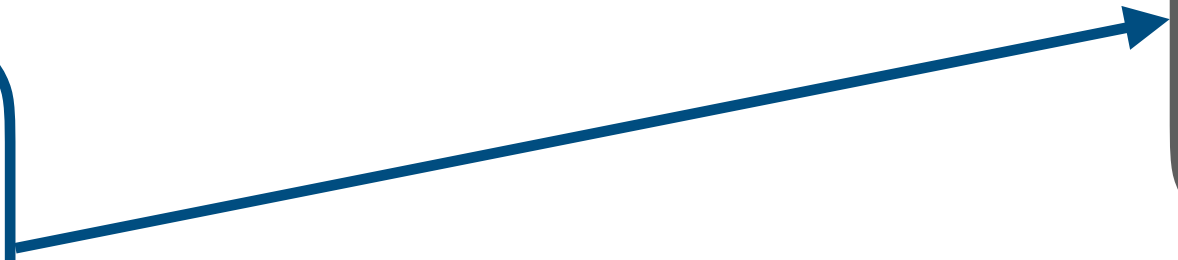
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a g b



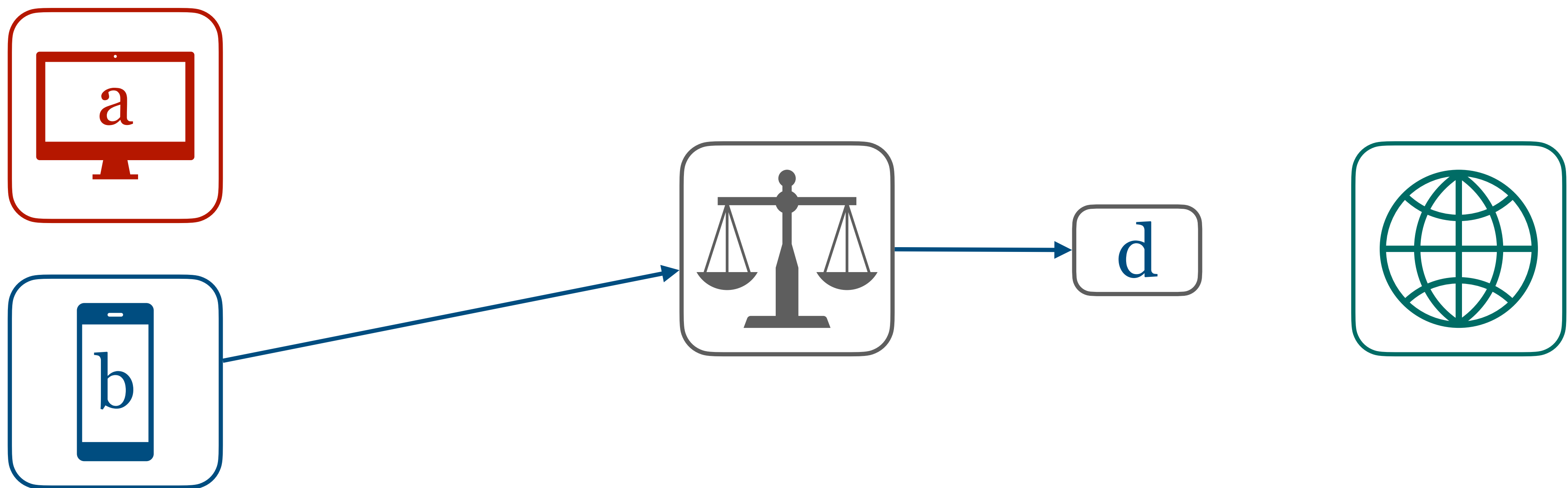


a g b d

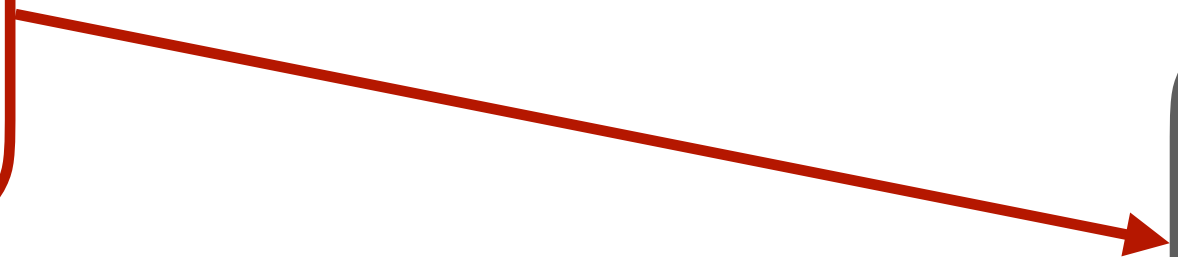


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a g b d b

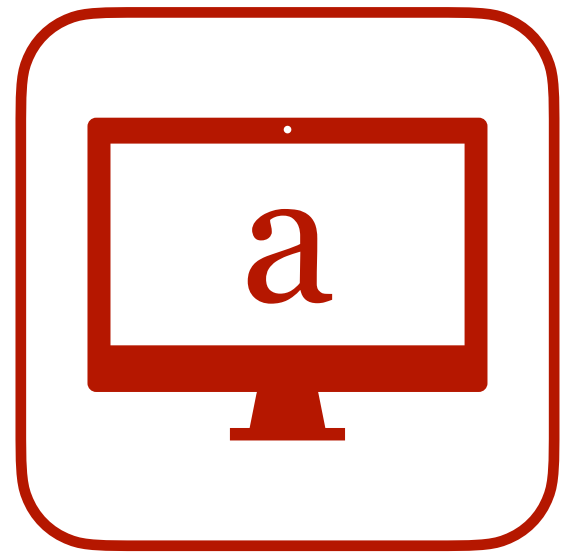


a g b d b d

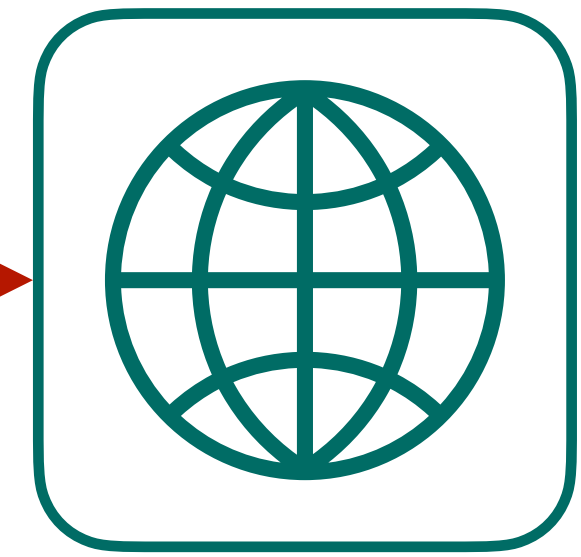


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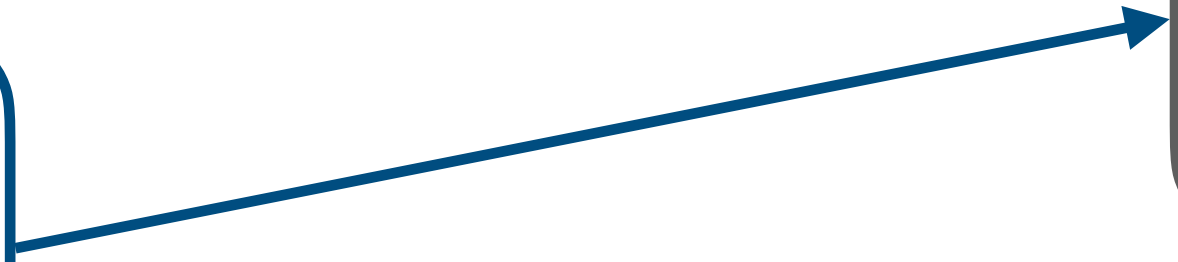
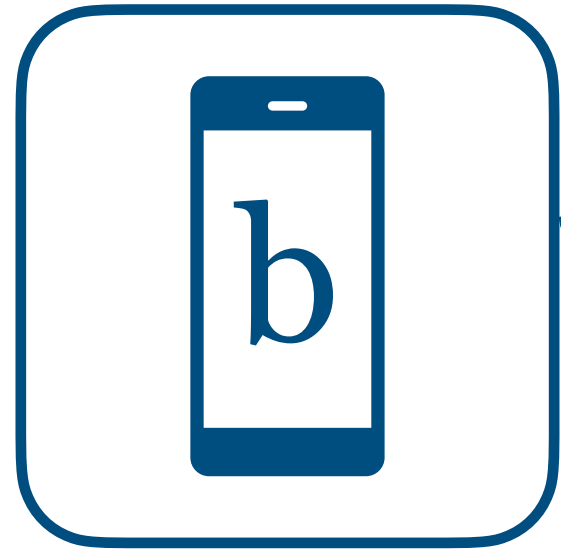
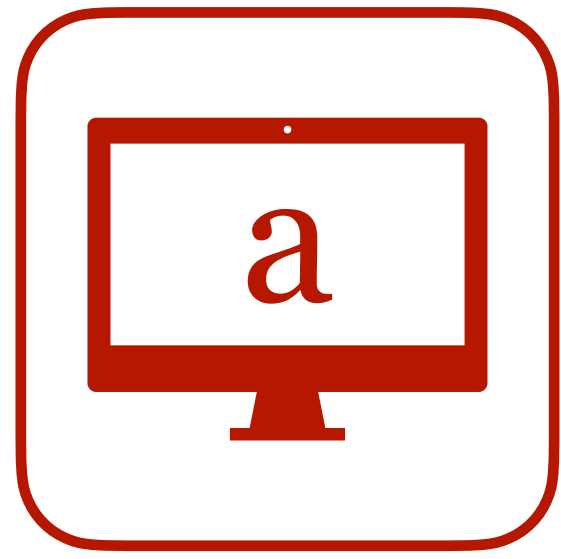
a g b d b d a



g

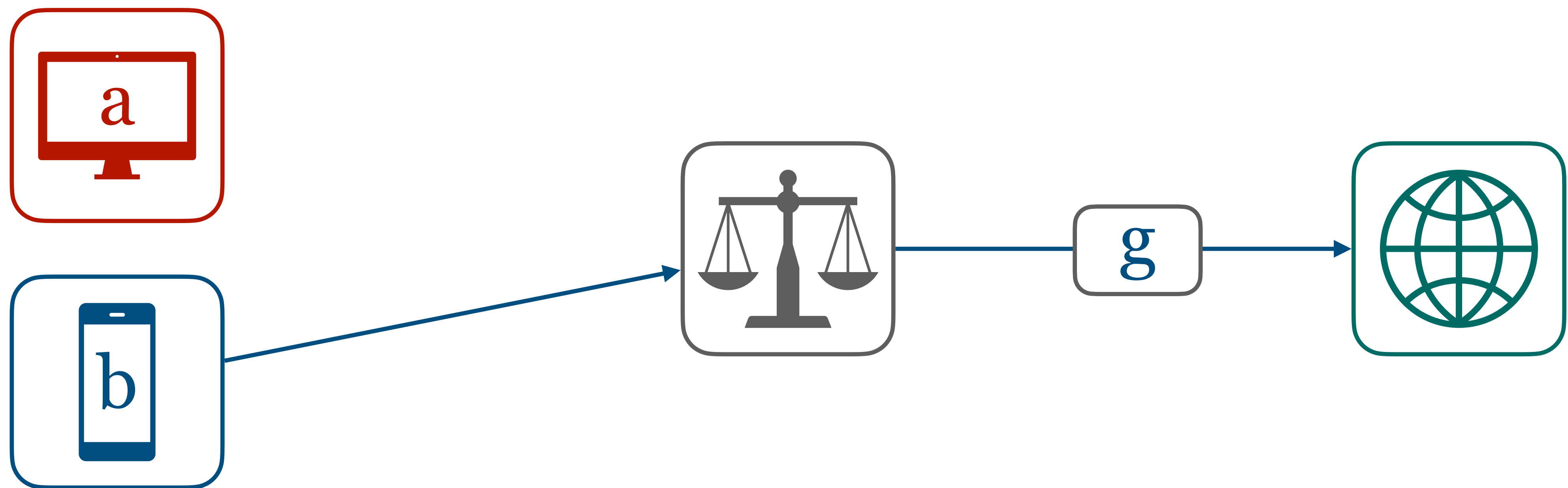


a g b d b d a g

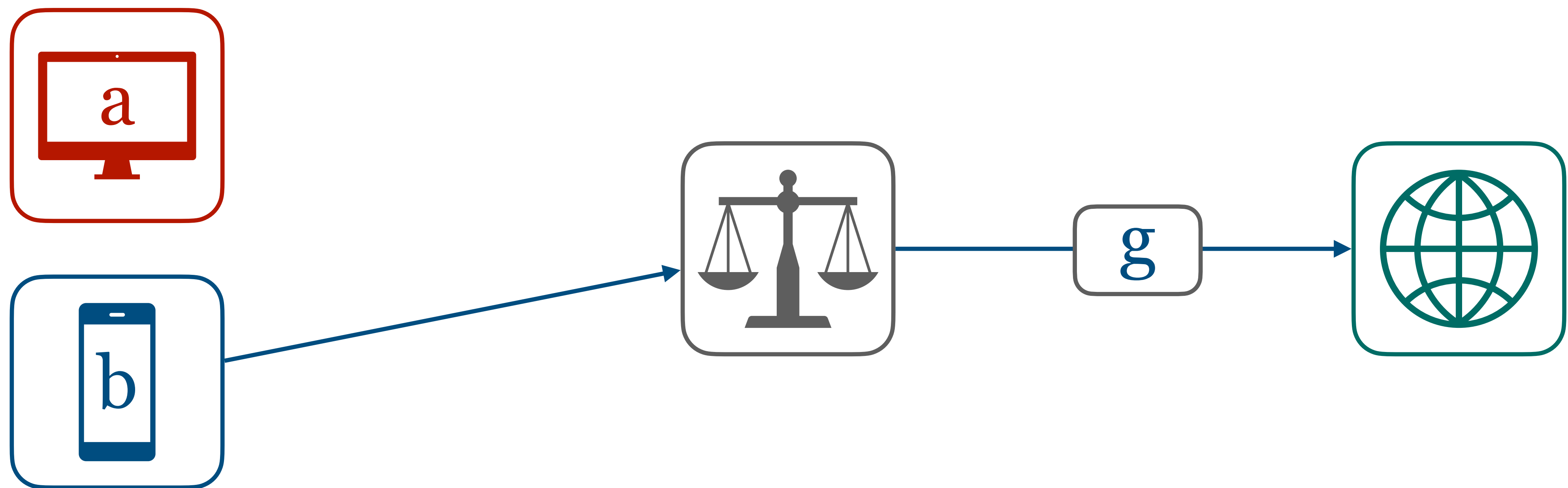


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a g b d b d a g b



a g b d b d a g b g



a g b d b d a g b g . . . . .



Property

$$\varphi \subseteq \Sigma^\omega$$

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a g b d b d a g b g . . . . .

Property

$$\varphi \subseteq \Sigma^\omega$$

Monitor

$$\mathcal{A} : \Sigma^* \rightarrow \{0, 1, ?\}$$

---

a g b d b d a g b g . . . . .

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Monitor

$$\mathcal{A} : \Sigma^* \rightarrow \{0, 1, ?\}$$

$$w \in \Sigma^\omega, u \prec w:$$

a g b d b d a g b g . . . . .

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$$\varphi \subseteq \Sigma^\omega$$

Monitor

$$\mathcal{A} : \Sigma^* \rightarrow \{0, 1, ?\}$$

$$w \in \Sigma^\omega, u < w:$$

$$\mathcal{A}(u) = 0 \Rightarrow w \notin \varphi$$

a g b d b d a g b g . . . . .

Property

$$\varphi \subseteq \Sigma^\omega$$

Monitor

$$\mathcal{A} : \Sigma^* \rightarrow \{0, 1, ?\}$$

$$w \in \Sigma^\omega, u < w:$$

$$\mathcal{A}(u) = 0 \Rightarrow w \notin \varphi$$

$$\mathcal{A}(u) = 1 \Rightarrow w \in \varphi$$

a g b d b d a g b g . . . . .

# Monitorability.

*If every infinite string has a point,  
where the monitor can stop watching.*

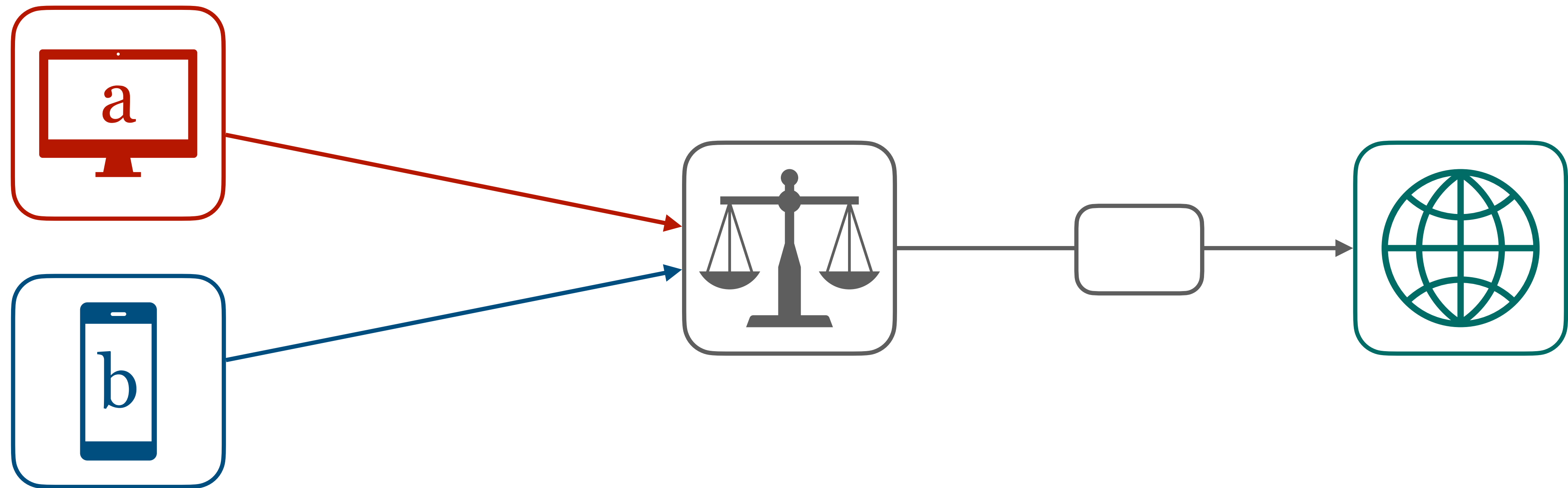
# Fairness Properties

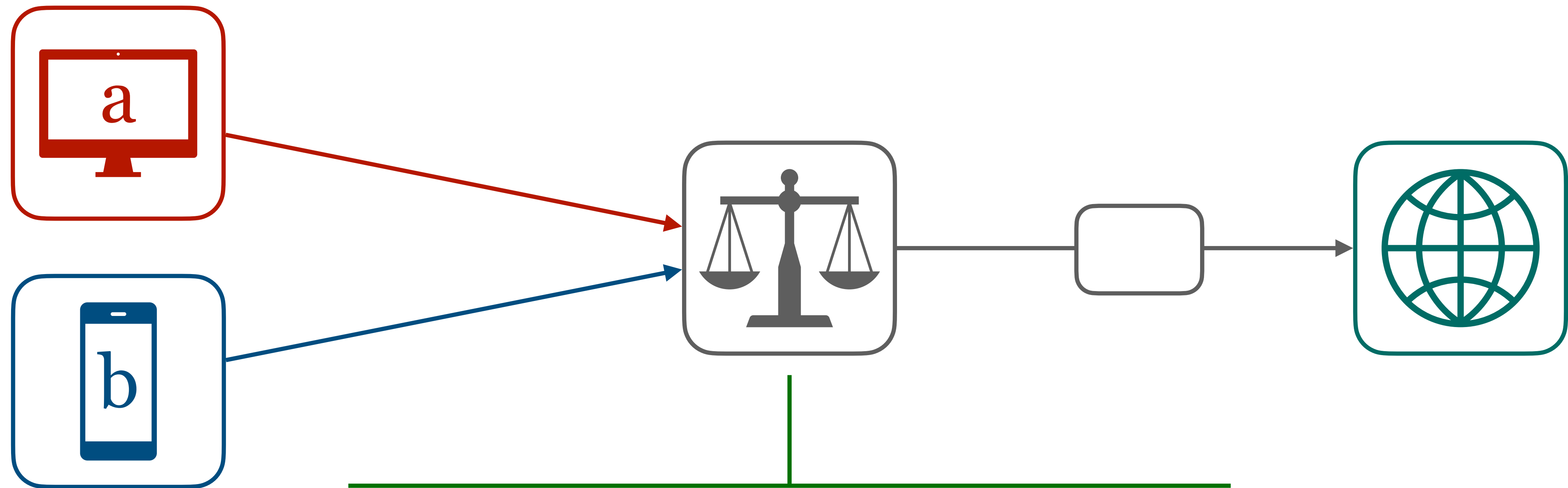
*From formal methods to machine learning.*

$$(\Box \Diamond a \rightarrow \Box \Diamond g) \wedge (\Box \Diamond b \rightarrow \Box \Diamond g)$$

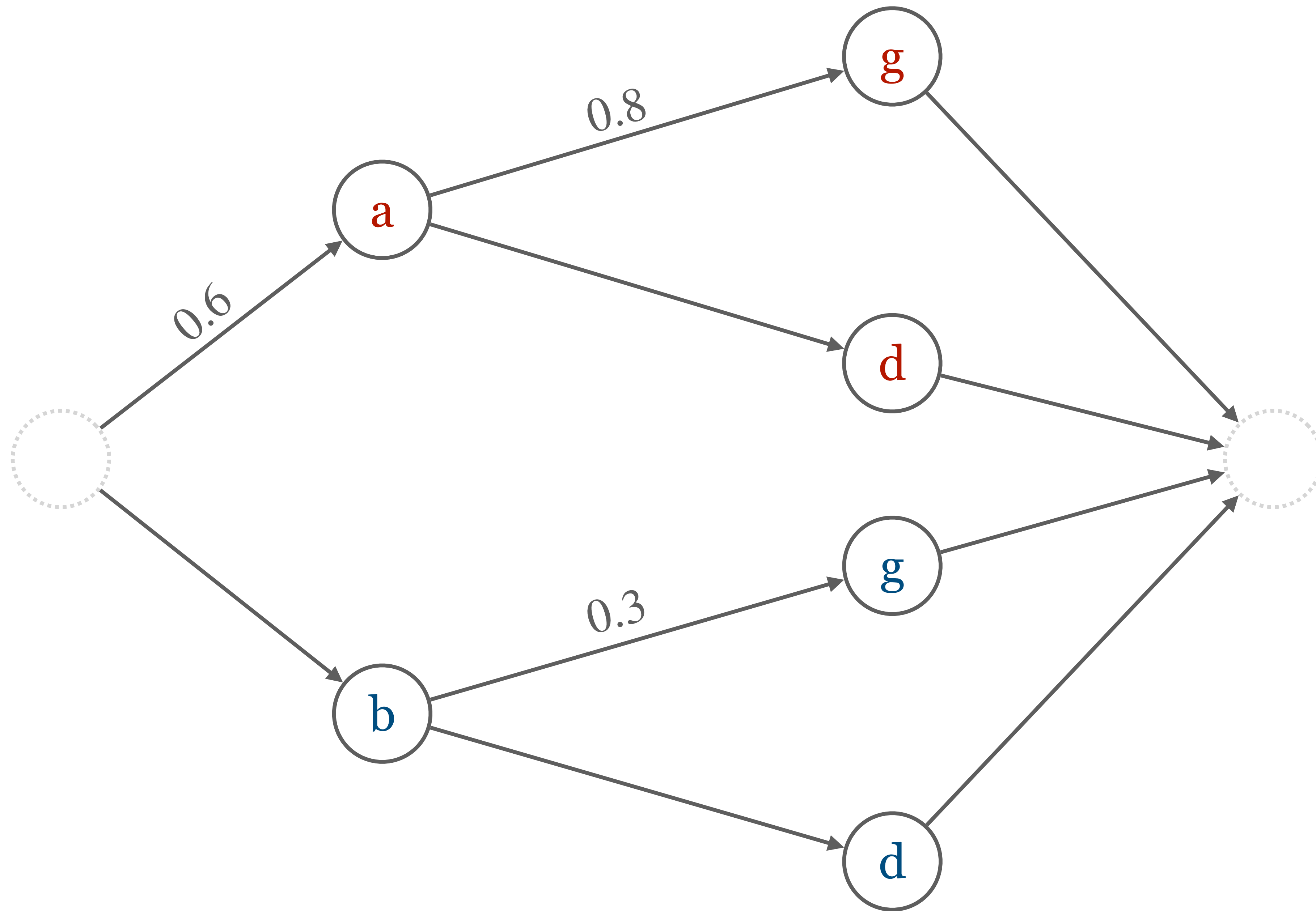
Fairness in Formal Methods







Memoryless & Probabilistic



$$(\Box \Diamond a \rightarrow \Box \Diamond g) \wedge (\Box \Diamond b \rightarrow \Box \Diamond g)$$

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$$\mathbb{P}(g \mid a) > 0 \wedge \mathbb{P}(g \mid b) > 0$$

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$$\mathbb{P}(g \mid a) > 0 \wedge \mathbb{P}(g \mid b) > 0$$



$$\mathbb{P}(g \mid a) - \mathbb{P}(g \mid b)$$

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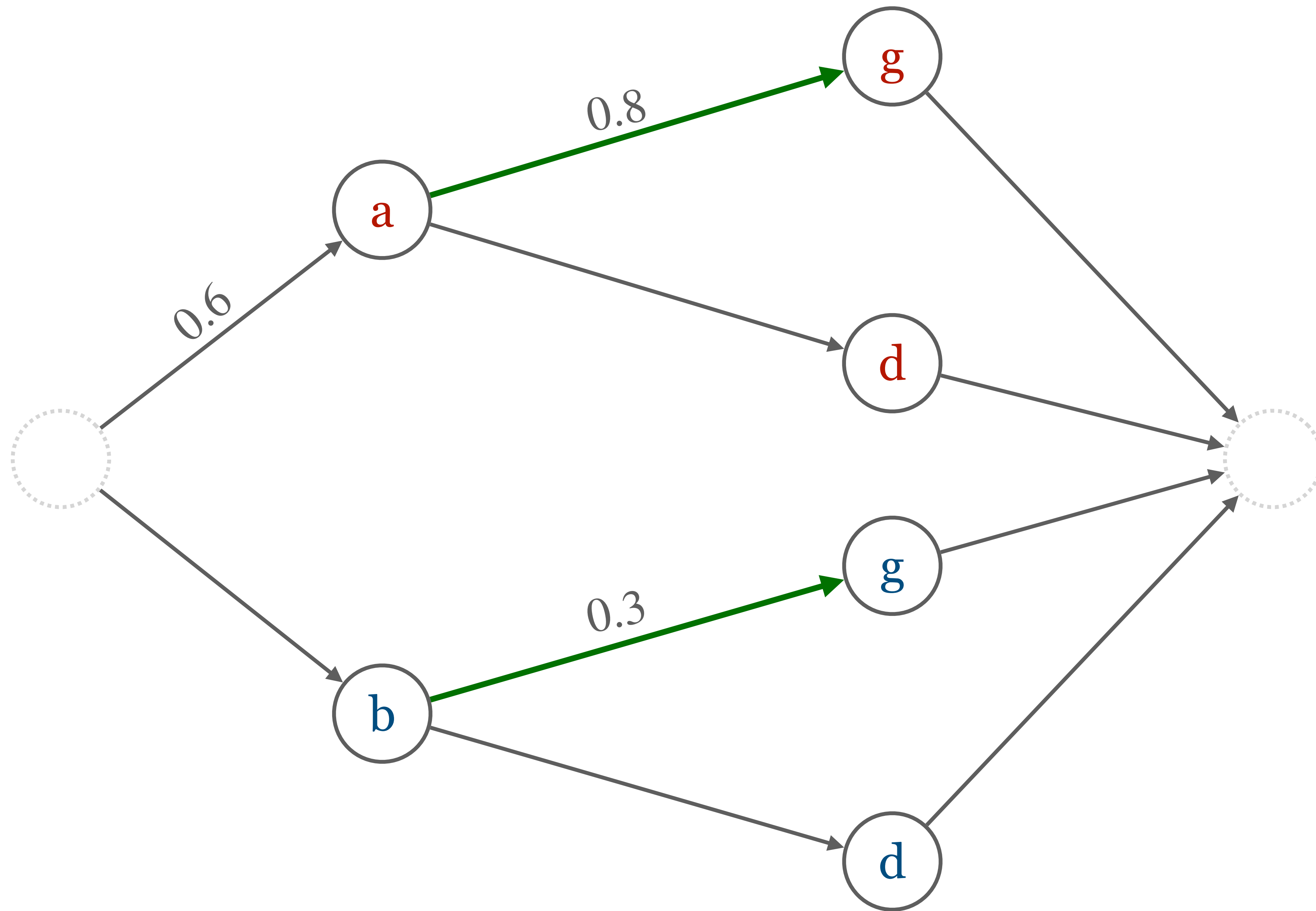
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Fairness in Machine Learning

# Monitorability.

*Can we estimate the property  
from observations?*

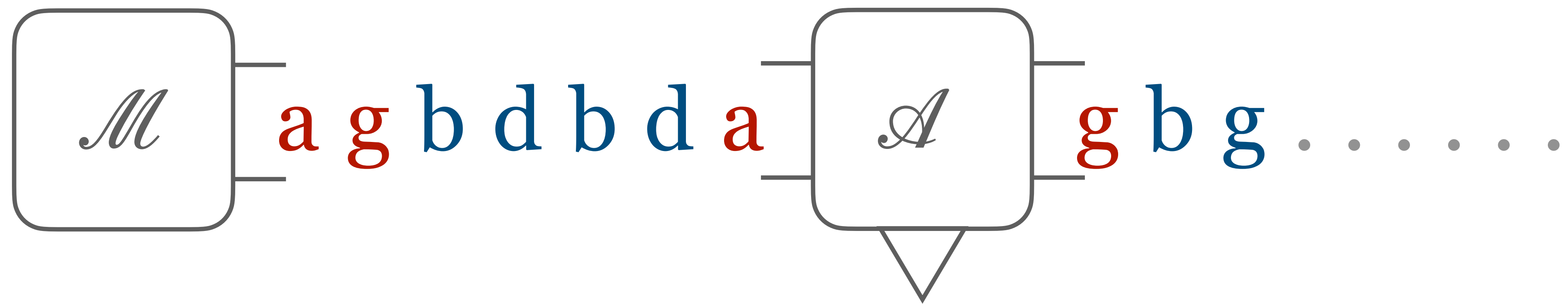


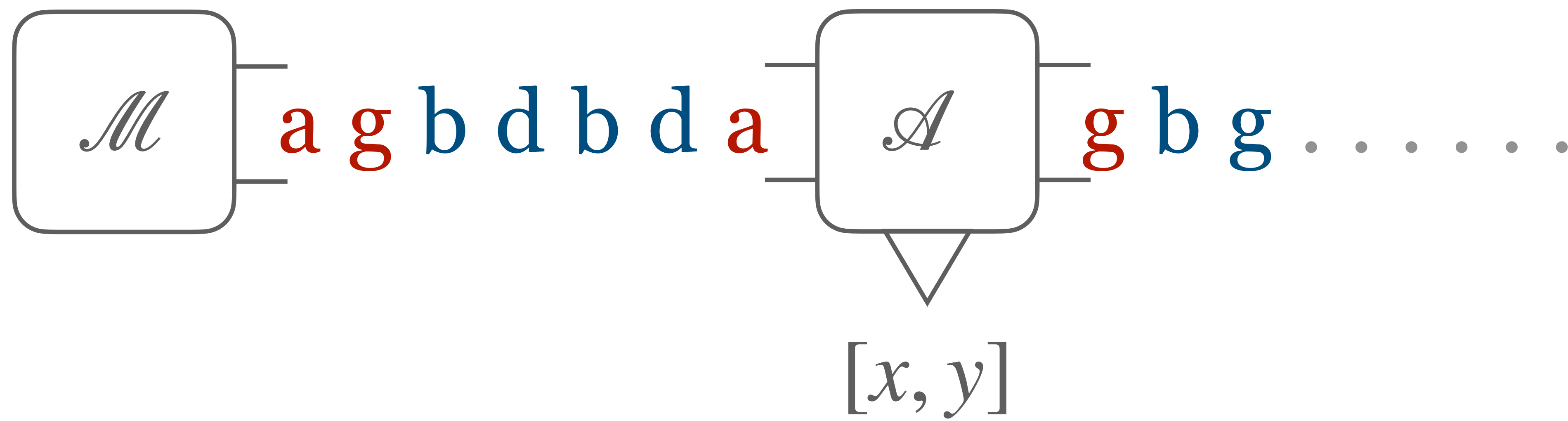


# General Idea.

*We observe a Markov chain and at every time step the monitor provides PAC-style guarantees.*

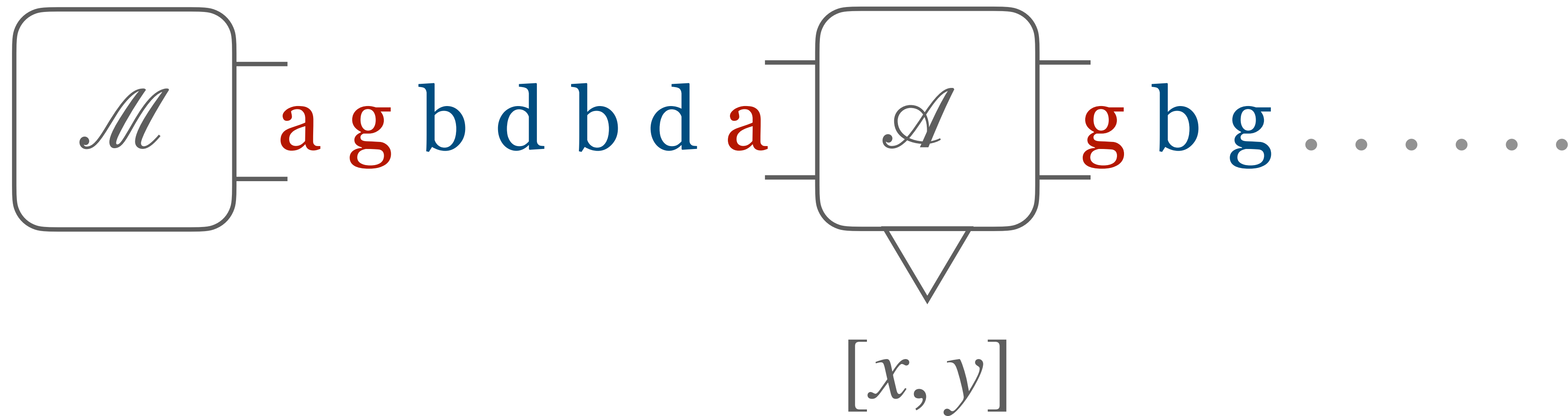






$$\mathbb{P}(g \mid a) - \mathbb{P}(g \mid b) \in [x, y] \text{ with probability } 1 - \delta$$


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# Problem Statement

*Let's be slightly more general.*

Let  $M \in \Delta(N-1)^N$ ,



Let  $M \in \Delta(N-1)^N$ ,  $W \sim (M, q_0)$

Let  $M \in \Delta(N - 1)^N$ ,  $W \sim (M, q_0)$  and  $U \prec W$ .

Let  $M \in \Delta(N-1)^N$ ,  $W \sim (M, q_0)$  and  $U \prec W$ . Given a function  $f: \Delta(N-1)^N \rightarrow \mathbb{R}$ ,

Let  $M \in \Delta(N-1)^N$ ,  $W \sim (M, q_0)$  and  $U < W$ . Given a function  $f: \Delta(N-1)^N \rightarrow \mathbb{R}$ , find a monitor  $\mathcal{A}: [N]^* \rightarrow \mathbb{R}^2$  such that:

Let  $M \in \Delta(N-1)^N$ ,  $W \sim (M, q_0)$  and  $U \prec W$ . Given a function  $f: \Delta(N-1)^N \rightarrow \mathbb{R}$ , find a monitor  $\mathcal{A}: [N]^* \rightarrow \mathbb{R}^2$  such that:

$$\mathbb{P}(f(M) \in \mathcal{A}(U)) \geq 1 - \delta$$

(Obviously we want the bounds to be as tight as possible.)

# Tradeoffs.

*We want to map the problem  
across four dimensions*

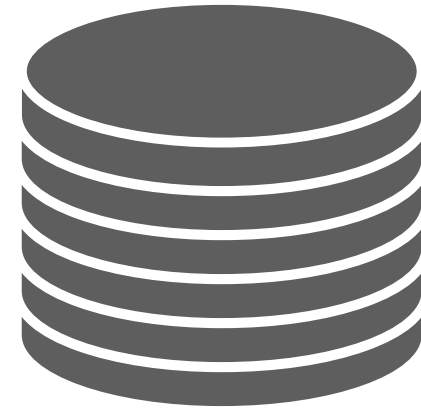
## Class of Functions



Class of Functions



Resource Complexity

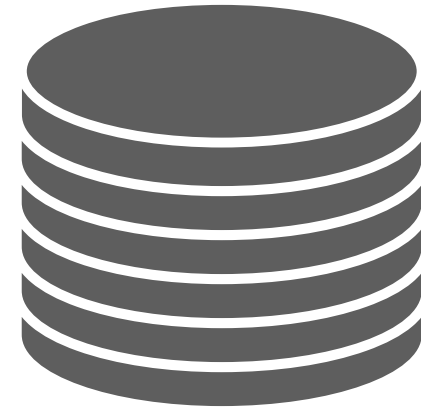




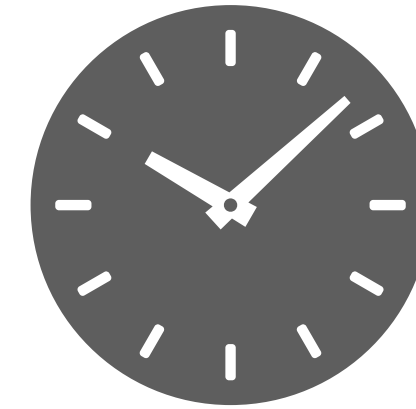
Class of Functions



Resource Complexity



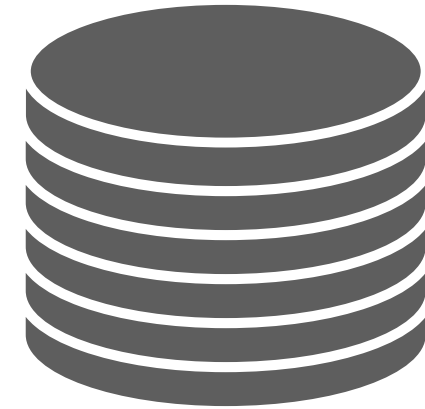
Time Complexity



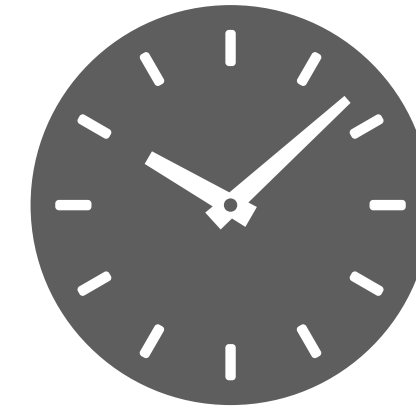
Class of Functions



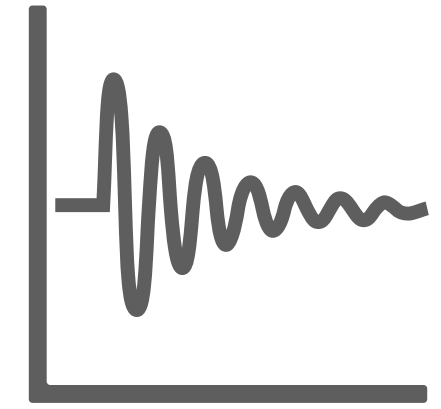
Resource Complexity



Time Complexity



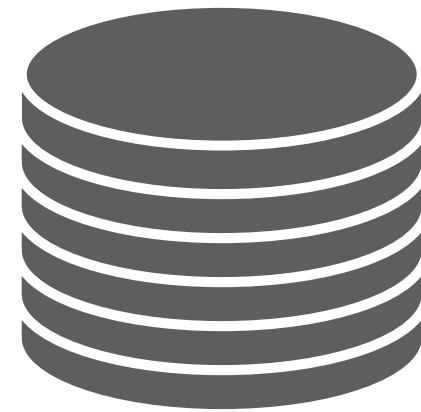
Sample Complexity



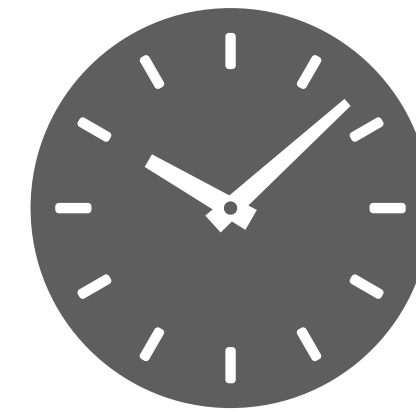
Class of Functions



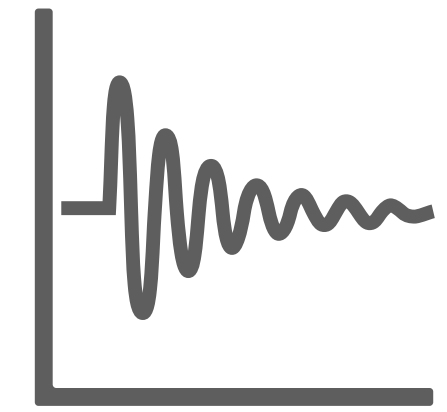
Resource Complexity



Time Complexity



Sample Complexity

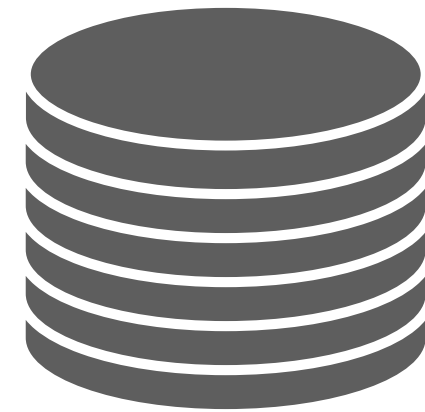


How does the class of functions influence the complexities, e.g. (in)dependent sums over  $M$ , polynomials over  $M$  (and/or the eigenvector of  $M$ ).

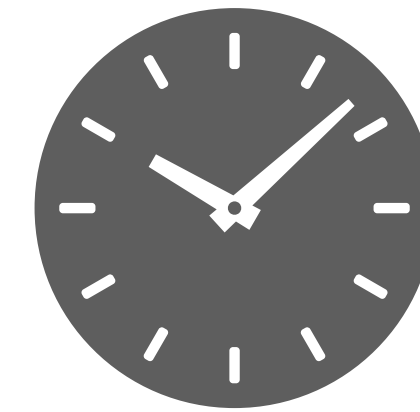
Class of Functions



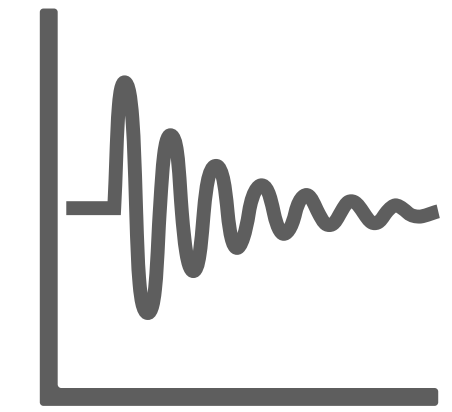
Resource Complexity



Time Complexity



Sample Complexity

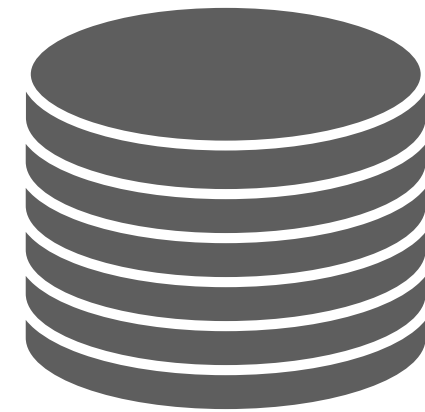


What is the minimal number of registers?  
(w.r.t. time/sample complexity)

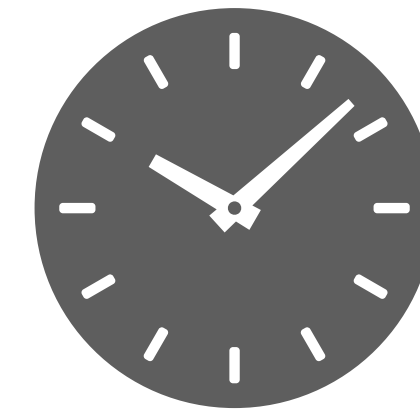
Class of Functions



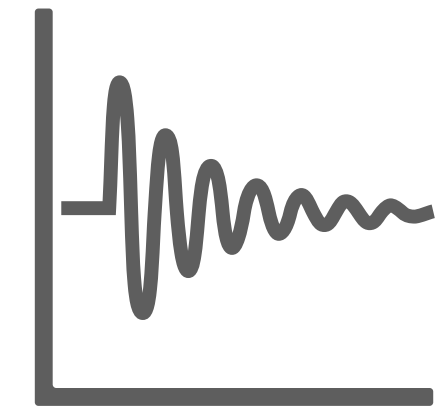
Resource Complexity



Time Complexity



Sample Complexity

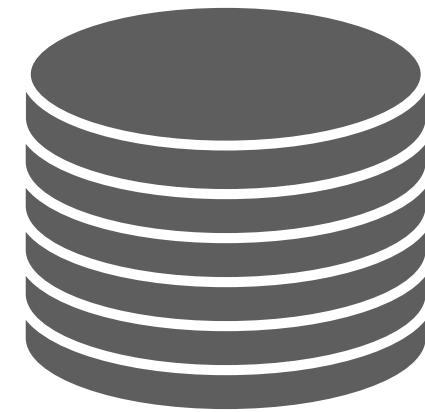


What is the minimal computation time?  
(w.r.t. resource/sample complexity)

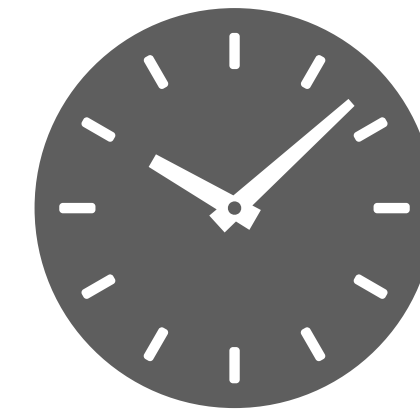
Class of Functions



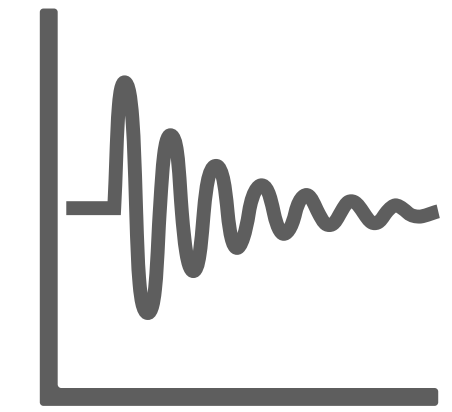
Resource Complexity



Time Complexity



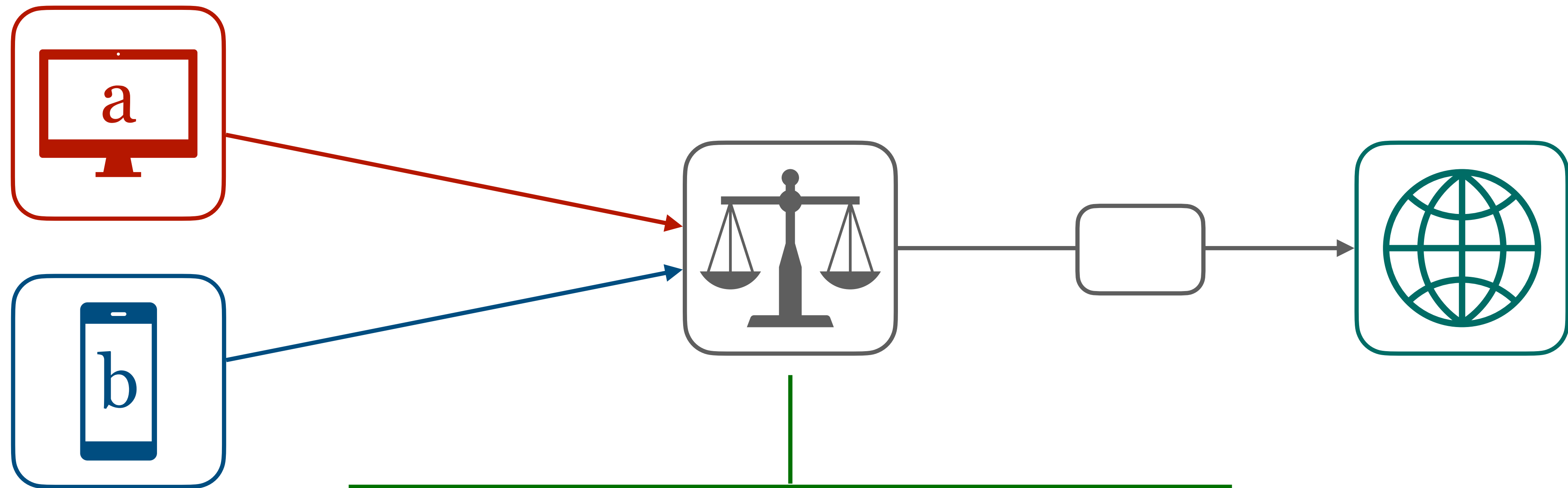
Sample Complexity



What is the rate at which the interval shrinks?  
(w.r.t. resource/time complexity)

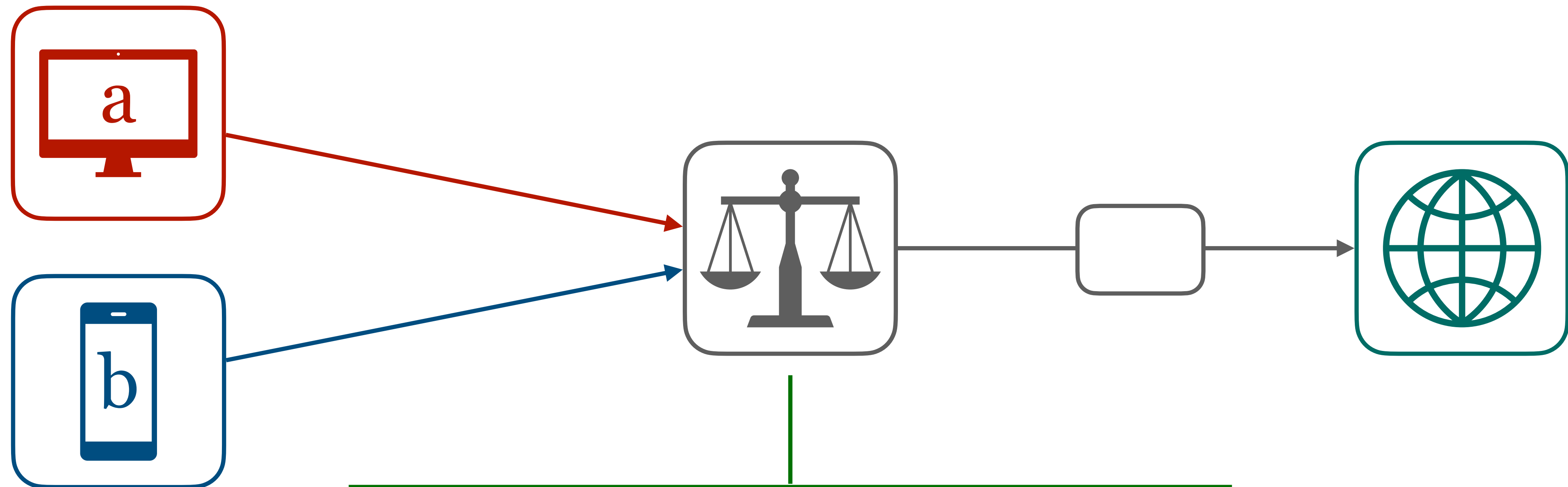
# But wait, there is more.

*What if the system is more complex?*



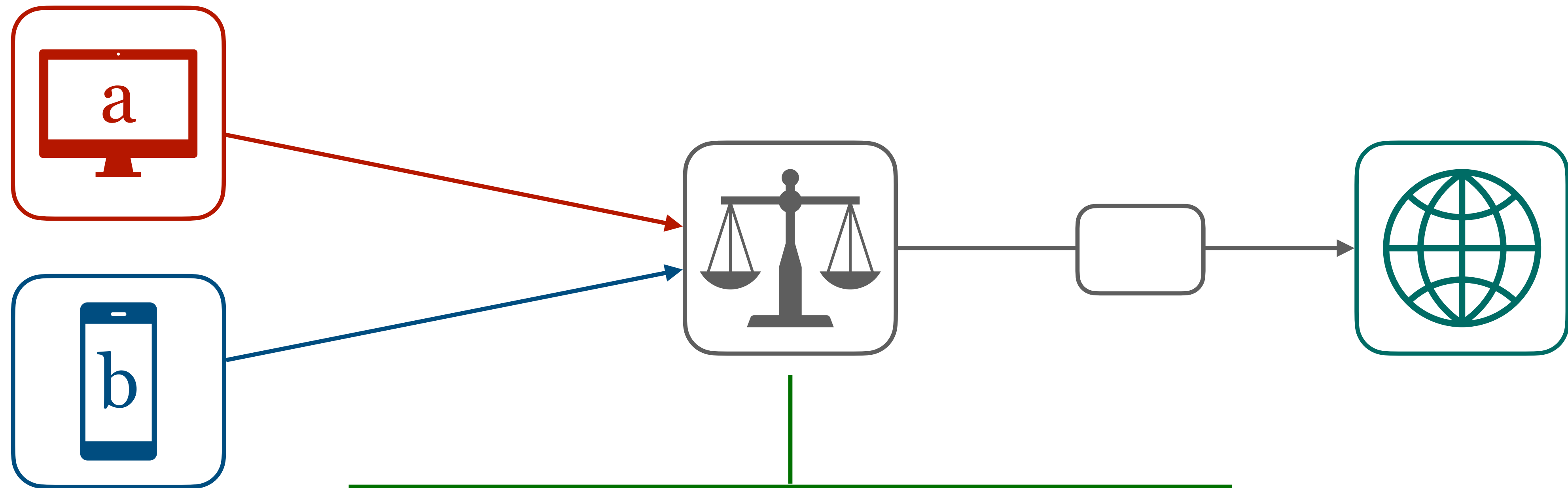
Remembers the last k-decisions.



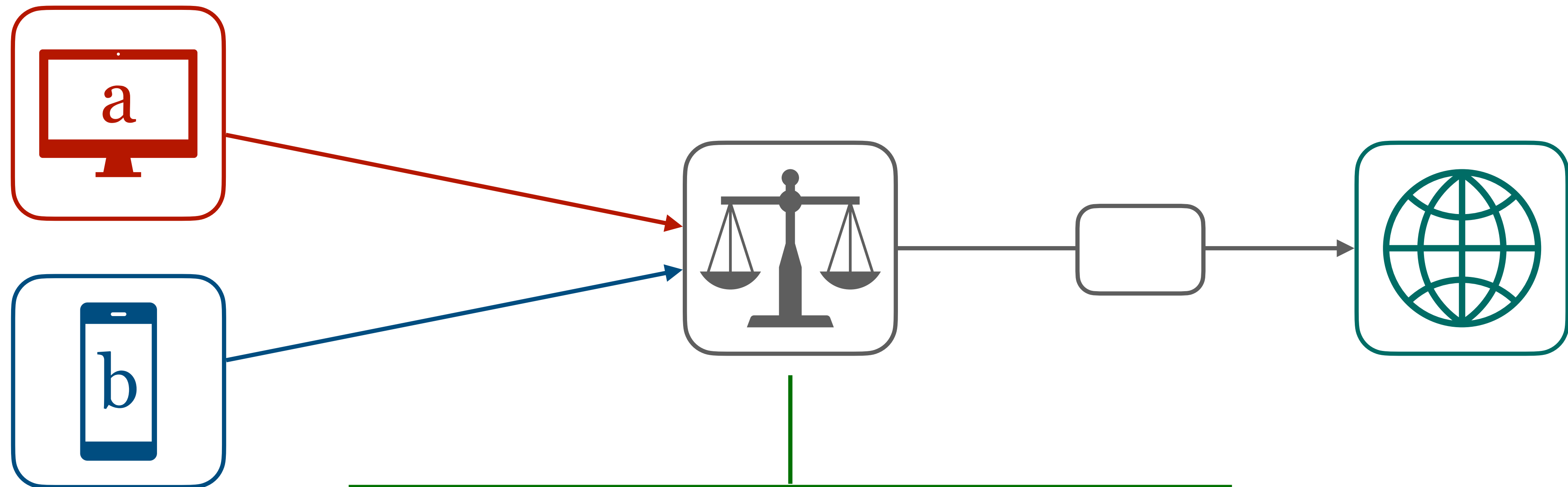


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Remembers some arbitrary  
k-decisions.



The policy changes at each time step,  
in a deterministic or probabilistic manner.



The decisions are corrupted  
or partially hidden.

# And many more.

*From our perspective there is a lot we don't know.  
It seems closely related to the concentration of  
functions over random variables  
with various dependencies.*

# What we did ...

*... so far.*

(Almost) arbitrary *arithmetic expressions* over  
*transition probabilities* of Markov chains.

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(Almost) arbitrary *arithmetic expressions* over *transition probabilities* of Markov chains.

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Efficient computation of *expectation* of arbitrary *polynomials* over *transition probabilities* of Markov chains in a *Bayesian* setting using a Dirichlet prior.

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(Almost) arbitrary *arithmetic expressions* over *transition probabilities* of Markov chains.

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Efficient computation of *expectation* of arbitrary *polynomials* over *transition probabilities* of Markov chains in a *Bayesian* setting using a Dirichlet prior.

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*Weighted sums* over *transition probabilities* of *time-inhomogeneous* Markov chains with *linear* and *observed change* in transition probabilities.



# Is this interesting to you?

*Let us know! (^\_^)*